ISSN 2395-1621

Double Guard: Detecting intrusions in Multitier web applications and providing security for front end and back end

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ABSTRACT

Nowadays internet services and applications have been widely used in daily life, allowing communication and personal information management from anywhere. To accommodate this increased use and complexity of data, web services have moved to a multilevel design, while the web server logic front-end application is executed and the data is outsourced to a server database or file. In this work double Guard, an IDS system that models the behavior of network user sessions on the server front-end Web and database back-end data is presented. By monitoring both web applications and subsequent databases, you may discover that IDS separate attacks would not be able to identify. Moreover, quantifying the limitations of any IDS multilevel based training sessions coverage and functionality. Double Guard is implemented using an Apache web server with MySQL and lightweight virtualization. Finally, the database is protected against manipulation of databases.

ARTICLE INFO

Article History

Received :16th April 2016

Received in revised form :

19th April 2016

Accepted : 21st April 2016

Published online :

25th April 2016

Keywords— Anomaly detection, virtualization, multitier web application

I. INTRODUCTION

Web services are being delivered and applications have been increased in popularity and complexity in recent years. Everyday tasks such as banking, travel and social networks, are all done through the web. These services usually employ a web interface server running the user interface logic of the application and a server back end consisting of a server database or file. Because of its ubiquitous use of personal corporate data and / or web services they have always been targeted. These attacks have become more diverse, as the focus has shifted to attack the front of the exploitation of the vulnerabilities of web applications in order to corrupt the system database back-end database (eg attacks SQL injection). A plethora of intrusion detection systems (IDS) are currently examining network packets individually both the web server and database system. However, there is very little work done on several levels anomaly detection (AD) systems that generate models of network behavior for both web and databases of network interactions. In this type of multi-level architecture, the database server back-end data often protected behind a firewall, while the web servers are remotely accessible via the Internet. Unfortunately, although they are protected from remote direct attacks, back-end systems are susceptible to attacks using web requests as a means of exploiting the rear end. To protect Web services from various levels, intrusion detection systems have been widely used to detect known attacks by matching patterns or signatures bad traffic use. A class of IDS fail of machine learning can also detect unknown attacks by identifying abnormal traffic network that deviates from the behavior called "normal" previously outlined during the IDS training. Individually, the Web IDS database and IDS can detect abnormal network traffic sent to any of them.However,we found that these IDS can not detect cases where normal traffic is used to attack the web server and database server.For example, if an attacker with administrator privileges can not log on to a Web server using the access

credentials for normal subscribers, he / she can find a way to issue a database query privileged data by exploiting vulnerabilities in the Web server. Neither IDS web or databases IDS would detect this type of attack because the Web IDS would only see traffic start typical user session and database IDS would only normal traffic of a user with privileges. This type of attacks can be easily detected if the IDS databases can identify a request for privilege from the Web server it is not associated with the privileged user access.Unfortunately,in the current architecture of the multithreaded web server, it is not feasible to detect or profile of this type of mapping causal link between the web server traffic and server database traffic and traffic cannot be clearly attributed user sessions.

Introducing Double Guard, a system used to detect attacks on web services at various levels. Our approach can model normal user sessions that include both isolated web interface (HTTP) and background (File or SQL) network transactions. To achieve this,we employ a technique lightweight virtualization to assign each user web session dedicated to a container, an isolated virtual computing environment. We use the container ID to associate the Web application accurately with the following database queries. Therefore, the double Guard can build a profile mapping causes taking into account the web server and DB traffic into account.

II. RELATED WORK

Double guard and its classification:-

Double Guard is a system used to detect attacks on web services of various levels [1] [2] System Intrusion Detection. A network can be classified into two types: anomaly detection and misuse detection. Anomaly detection IDS requires first to define and characterize the proper and acceptable static and dynamic behavior of the system, which can then be used to detect abnormal changes or abnormal behavior. The boundary between acceptable and abnormal forms of data stored code and is precisely definable. Behavioral models are constructed by performing a statistical analysis of historical data or by using rule-based to specify the behavior patterns approaches. An anomaly detector then compares the actual usage patterns against established models to identify abnormal events [1] [2] [3].

III.METHODOLOGY

This approach can model normal user sessions that include both isolated web interface (HTTP) and background (File or SQL) network transactions [1] [3]. a virtualization technique lightweight [1] is used [3] to assign each session users to a dedicated container, an isolated virtual computing environment. the container ID is used to associate the Web application accurately with the following database queries. Guard double container-based IDS form with the entry of several sequences to produce alerts. The correlation of inflows provides a better characterization of the system for detecting anomalies because the intrusion sensor has a model of normality more accurate to detect a wider range of threats [1] [2] [3] [4].

Possible Attack:-

Some of the major attacks are generally used by attackers to piracy i.e. SQL injection, Direct DB Attack, Hijack future session attack, Privilege escalation [1] [2] [3] [5] and D-DOS attack [1].

Algorithm Used:-

In order to detect such attacks used algorithms are static model building algorithm models [1] [2] [3] [4].

Limitations:-

Vulnerabilities Due to Improper Input Processing:-

Once the user enter malicious are normalized, Double Guard can not detect hidden attacks on the values [1].

Possibility Of Evading Double Guard:-

It is possible for an attacker to discover mapping patterns by doing the code analysis or reverse engineering, and issue expected web requests prior to performing malicious database queries [1].

Distributed DOS attacks:-

Previous double Guard system was not designed to mitigate attacks like D-DOS. Such attacks can also occur in server architecture without the database back-end database. denial of service attacks are common and very fashionable these days. In denial of service attack, the attacker attempts to prevent legitimate users from using a service or the closure of a service because some vulnerability to crash the machine application [1].

IV. DOUBLE GUARD SYSTEM ARCHITECTURE

SYSTEM ARCHITECTURE:-

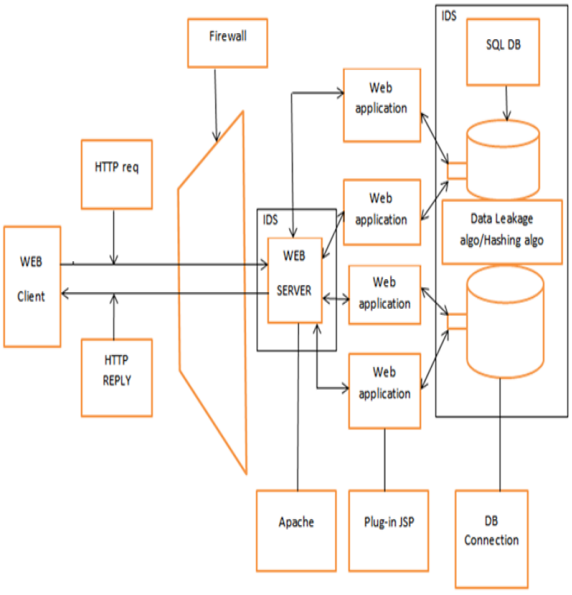

Figure 1: System Architecture diagram

This is the proposed system architecture design. In this, first the client send request for information related to particular product and product price then request is identify if the request

is HTTP request or a query and this is done using static model building algorithm. After the request is categorized if the request is HTTP request then that request is passed through firewall and web server receives that request and that request is send as a query to database server and response is sent accordingly but the request is handled only after user authentication is satisfied. If the value in database is changed then data leakage occurs and using data leakage algorithm we save the record from attackers and report back to admin. If user request for data which is hack then previous data is provided to that user using hashing algorithm.

Attacks scenario:-

1. SQL-Injection Attacks:-

Attacks such as SQL injection do not require compromise the web server. Attackers can use the existing vulnerabilities in the web server logic inorder to inject the contents of data or a string containing the exploits and then use the web server to relay these exploits to attack the database back-end database. Since our approach provides detection of two levels, even if exploits are accepted by the web server, the content transmitted to the server database would not be able to take the envisaged structure for the application of certain web server. Attacks such as SQL injection do not require compromise the web server. Attackers can use an existing vulnerabilities in the web server logic inorder to inject the contents of data or a string containing the exploits and then uses the web server to relay these exploits to attack the database back-end database. Because our system allows detection of two levels, even if exploits are accepted by the web server, the content transmitted to the server database would not be able to take the envisaged structure for the application of certain web server.

For example, SQL injection attack changes the structure of SQL query or queries, even if the data injected where to go on the side of the web server, SQL queries are generated in a different structure that could be detected as a deviation from the structure SQL query that usually follow such web application

2. D-Dos Attacks:-

Double guard is not designed to mitigate denial of service attacks D- [Figure 3]. These attacks can also occur in server architecture without the backend data base. In computing, a service denial (DoS) is an attempt to make an appeal of the machine or network is unavailable to its intended users, to interrupt or suspend the services of a host connected to the Internet temporarily or indefinitely . A distributed denial of service (DoS-D) is where the source of attack is more than one, and often thousands-unique IP address.Criminals are responsible for DoS attacks that often target sites or services hosted on high-profile web servers, such as banks, credit card payment gateways; but the motives of revenge, blackmail or activism may be behind other attacks. A D-DoS attack i.e, distributed denial of service attack that occurs when multiple systems flood the bandwidth or resources of targeted system, usually one or more web servers. Such an attack is usually the result of multiple compromised systems that flood the target system with traffic.

When a server is overloaded with connections, new connections can no longer be accepted. The main advantages of an attacker using a distributed denial of service that multiple machines can generate more attack traffic than one machine, several machines are more difficult to attack the machine off an attack,and that the behavior of each machine can be stealthier attack, so it is harder to track and off. These advantages make the attacker challenges for defense mechanisms. For example, buying more incoming bandwidth is limited to the current volume of the attack could not help, because the attacker might be able to simply add more machines to attack. This, after all will end completely crashing a website for periods of time. The malware mechanisms can lead to denial-of-service attack D-; one of the best known examples of this was my destiny. DoS mechanism was triggered at a specified date and time. This type of D-DoS involved hard coding the IP destination address before the release of malware without additional interaction was needed to launch the attack.

3. Direct DB attack:-

It is possible for the attacker to bypass the webserver or firewalls and connect directly to database. An attacker could also have been already taken over the webserver and be submitting such queries from the webserver without sending the web requests. Without matched web requests for such queries, a webserver IDS could detect neither. Furthermore, if these database queries were made within the set of an allowed queries, then the database IDS would not be able to detect it either. However, this type of an attack can be caught with our approach since we cannot match any web requests with these queries.

V. RESULT

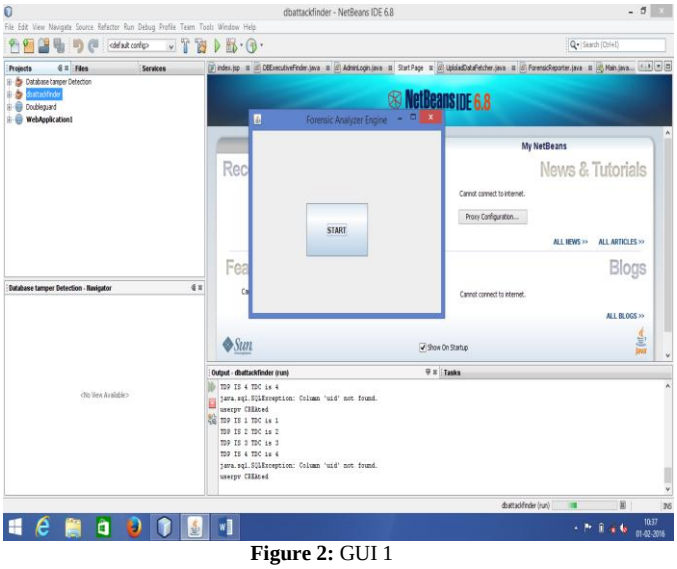

Figure 2: GUI 1

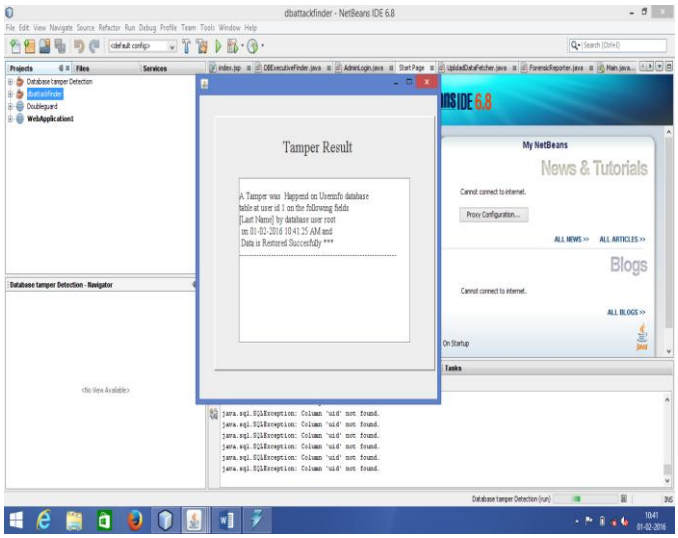

Figure 3. GUI 2

This proposed system consists of forensic analyzer engine as shown in figure 2, database tamper detection system, a database consisting all data, tamper result viewer as shown in figure 2. As soon as the forensic analyzer engine is started it goes to database tamper detection system which can be accessed after providing authenticated login details after which one can start engine and access database and hack and alter database. If database is altered then the tamper details are sent immediately to admin providing admin with details like which data is altered at what time and date also which root database is used to do so and also the altered data is restored to its original values so as to provide security.

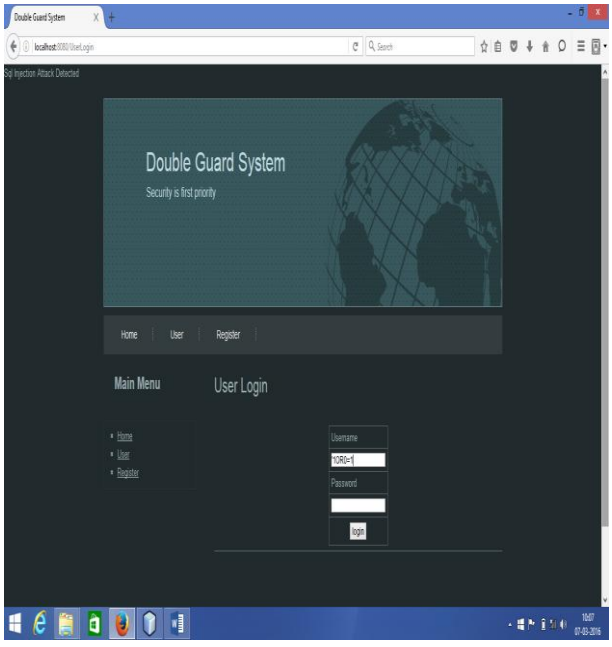

Figure 4: SQL-Injection Detection

Normally any systems where user needs to use this system register itself to use this system they fill information about itself. Attackers use SQL query to use system without registration. Our system provides security against SQL attacks detected and no access further.

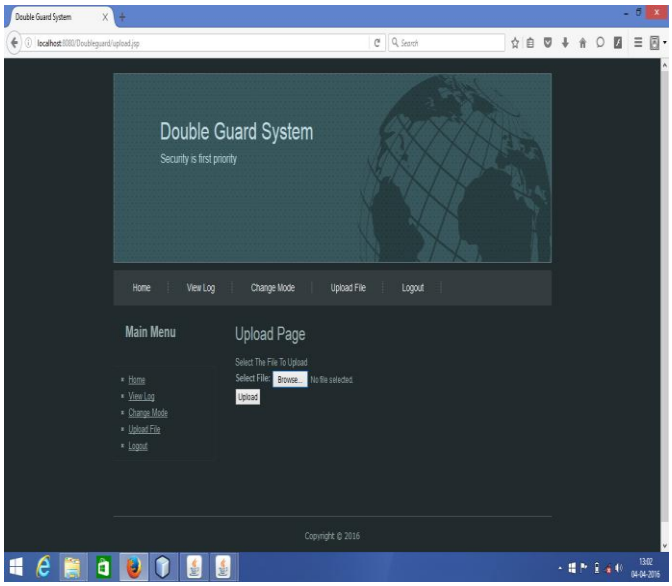

Figure 5. GUI 3

Whenever admin sends data as response the data is uploaded successfully but when that data exceeds size of 1MB it will not be uploaded due to D-DOS attack this is because if data requested is send by admin again and again for just one user other goes to waiting state so in order to avoid that this condition have been kept.

VI. ALGORITHM

Static Model Building Algorithm (I):-

Ensure: The Mapping Model for static website.

Input: Set AQ for database query. Set AR for server request.

Step 1: Identify the input type of HTTP request whether it is a query or a request.

Step 2: for each different request do, if r is a request to static file.

Step 3: Store inputs in hash table as per their type AQ for query and for request AR.

Step 4: The key for hash table entry will be set as the input itself.

Step 5: Forward AQ and AR to virtual server to validate.

Step 6: If attack identified then virtual system automatically terminate the HTTP request

Step 7: Else HTTP request is forwarded to the original server.

Step 8: Display information.

Step 9: Exit.

Data leakage algorithm (II):-

Input: Input data D = D1, D2, D3... Dn saves into the hash table.
Step 1: Arrange all input data into matrix format (save into log files).

Step 2: Consider m as a selected data act as a new selected data.

Step 3: m position gets changed after allocated time period.

Step 4: If Ms data get hacked.

Step 5: Data leakage is occurs.

Step 6: We have to check the leakage data and prevent

Step 7: Using Revert back function we have to get original data.

Step 8: When user calls that corrupted file, hash function gives to user a previous data.

Step 9: Return True.

MD5 Hashing algorithm (III):-

MD5 stands for Message Digest algorithm 5 which is a widely used cryptographic hash function.The idea behind this algorithm is to take up a random data (text or binary) as an input and generate the fixed size hash value as the output The input data can be of any size or length, but the output hash value size is always fixed

Step 1: Start

Step 2: For each candidate set element.

Step 3: For PV (i) and CV (i) compare attributes and detect which fields are corrupted.

Step 4: get who and when of corruption event.

Step 5: Prepare a report.

Step 6: Stop

VII. FUTURE SCOPE

The basic idea is provide two tier securities to for web applications. The aim is to secure the web server from the attacker client and to secure the data from the internal authorized persons in the data care centers. This security model can be further extended to provide security against other attacks.

VIII. CONCLUSION

Hereby we conclude that in this project we have successfully provided security for front end and back end i.e. preventing attacks like SQL injection and D-DOS attack at front end and Direct DB attack at back end using two intrusion detection system and totally secure the system.

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